

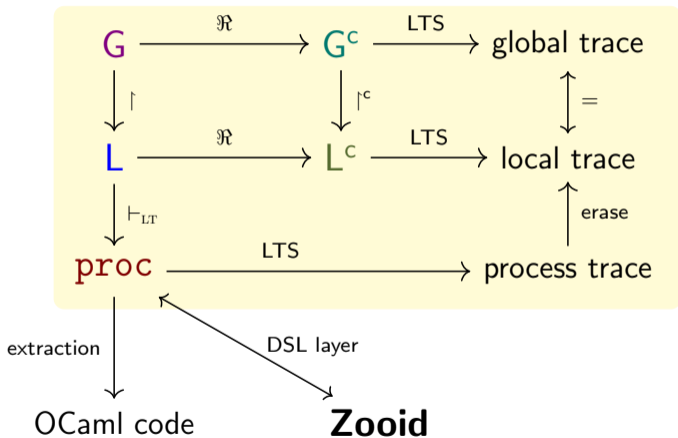
Act II

Smol-Zooid: multiparty with shallower embedding

Goals

1. Certifying **individual** processes of a **distributed system**
2. Extracting runnable code
3. Avoiding complex formalisations of binders, whenever possible

Overview



Smol Zooid

- We combine **shallow/deep embeddings** of binders
 - We use DeBruijn indices for the deeply embedded binders
- SZooid constructs are **well-typed by construction**
- We leverage **Coq code extraction** mechanism
- For simplicity, SZooid does not cover choices

Core Processes

In: <http://github.com/emtst/gentleAdventure>

```
Inductive proc :=
| Inact | Rec (e : proc) | Jump (X : nat)
| Send (p : participant) {T : type}
  (x : interp_type T) (k : proc)
| Recv (p : participant) {T : type}
  (k : interp_type T -> proc)

| ReadIO {T : type} (k : interp_type T -> proc)
| WriteIO {T : type} (x : interp_type T) (k : proc).
```

Payload Types

We need to define a type for payload types:

- We need a decidable equality on payload types
- We need a decidable equality on payload values

```
Inductive type := Nat | Bool | ...
```

```
Definition interp_type : type -> Type := ...
```

Semantics: events

The semantics is an LTS:

- the labels are the **communication events**
- it is parameterised by a **payload interpretation function**
- traces are obtained as the greatest fixpoint of the LTS step

Inductive action := a_send | a_recv.

Record event interp_payload :=

```
{ action_type   : action;
  subj          : participant;
  party         : participant;
  payload_type  : type;
  payload       : interp_payload payload_type
```

```
}
```

Semantics: Recursion Variables and I/O

`p_unroll` : `proc` \rightarrow `proc`

`p_unroll` exposes the first communication action in a process:

- “runs” any I/O action
- unfolds recursion

Definition `p_unroll` : `proc` \rightarrow `proc` := ...

Semantics: step

The step of the LTS is defined as a **function**:

```
Definition step' e E :=  
  match e with  
  | Send p T x k =>  
    if (action_type E == a_send) && (party E == p) &&  
      (eq_payload (payload E) x)  
    then Some k else None  
  | Recv p T k => ... | _ => None  
end.
```

```
Definition step e := step' (p_unroll e).
```

Local Types

We introduce a typing discipline that associates processes with **local types**, that characterise their communication behaviour:

```
Inductive lty :=  
  | l_end  
  | l_jump (X : nat)  
  | l_rec (k : lty)  
  | l_send (p : participant) (T : type) (l : lty)  
  | l_recv (p : participant) (T : type) (l : lty).
```

Type System

```
Inductive of_lty : proc -> lty -> Prop :=
| lt_Send      p T k L x :
  of_lty k L -> of_lty (@Send p T x k) (l_send p T L)
| lt_ReadIO    T k L :
  (forall x, of_lty (k x) L) -> of_lty (@ReadIO T k) L
| ...
.
```

Smol Zooid: Smart Constructors (I)

- It would be tedious to type up both a local type and a process
- Users would need to provide a proof that processes are well-typed

We define **SZooid** (Smol Zooid), to write well-typed processes by construction, avoiding repetition.

Smol Zoid: Smart Constructors (and II)

Definition `SZoid L := { p | of_lty p L}`.

Definition `z_Send p T x L (k : SZoid L)`
`: SZoid (l_send p T L)`
`:= exist _ _ (lt_Send p x (proj2_sig k)).`

...

Inferring Local Types

SZooid constructs fully determine their types from their inputs, so we can ask Coq to infer local types associated with SZooid terms:

Definition AZooid := { L & SZooid L }.

Subject Reduction

Theorem preservation (e : proc) (L : lty)
(H : of_lty e L) (E : rt_event) :
forall e',
step e E = Some e' ->
exists L', lstep L (ev_erase E) = Some L' /\
of_lty e' L'.

Extraction

- We convert `proc` to function calls in an **ambient monad**
- We extract the monadic code to OCaml
- The ambient monad needs to be implemented in OCaml
- Processes are extracted using Higher-Order modules, so it is straightforward to change the underlying transport
- **Remark:** SZooid does not provide an implementation of the ambient monad, but Zooid does, using TCP/IP sockets:
<https://github.com/emtst/zooid-cmpst>

Extraction Module

```
Module ProcExtraction (MP : ProcessMonad).  
  Fixpoint extract_proc (d : nat) (p : proc) : MP.t unit  
    := match p with  
       | Send p T x k  
         => MP.bind (MP.send T p x)  
                  (fun=> extract_proc d k)  
       ...  
End ProcExtraction.
```

Example Extraction

```
Module ALICE (MP : ProcessMonad) : PROCESS(MP).  
  Module PE := ProcExtraction(MP).  
  Definition proc :=  
    Eval compute in PE.extract_proc 0 alice.  
End ALICE.
```

Extraction ALICE.

Summary

- We have seen how to encode a small calculus of Multiparty Processes, with a basic type system
- **Next: how do we relate traces of individual processes to a larger system?**